

# Efficient Quantum-Safe Communication using Hybrid Encryption

Rei Safavi-Naini



UNIVERSITY OF  
CALGARY

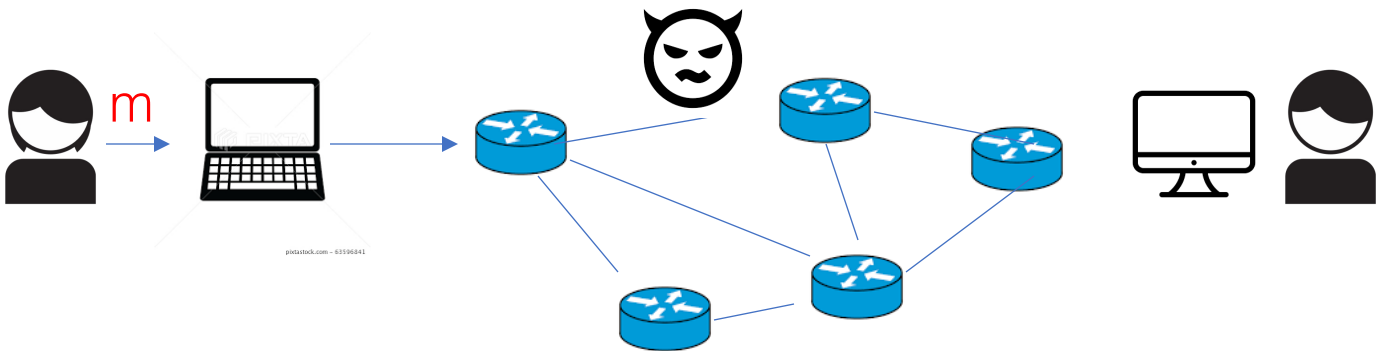
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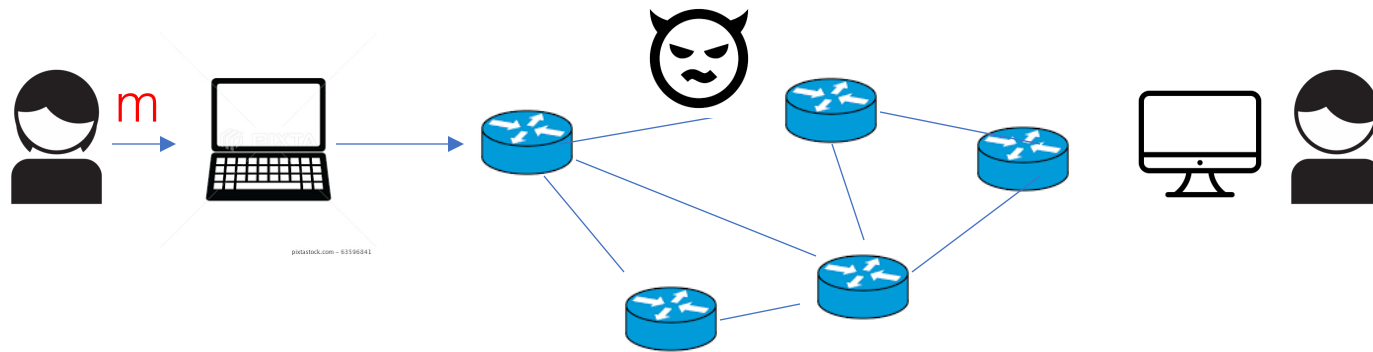
# Plan of the talk

- Establishing a confidential channel (storage)
  - Approaches
- Hybrid encryption (KEM/DEM)
  - Confidential channel with proved security
- A new approach to HE
  - Correlated randomness model
- Combing KEMs
- Concluding

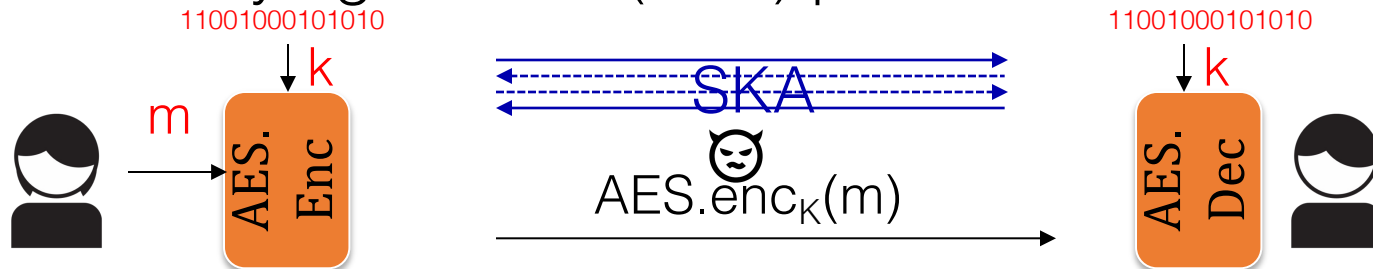
# Confidential communication over Internet



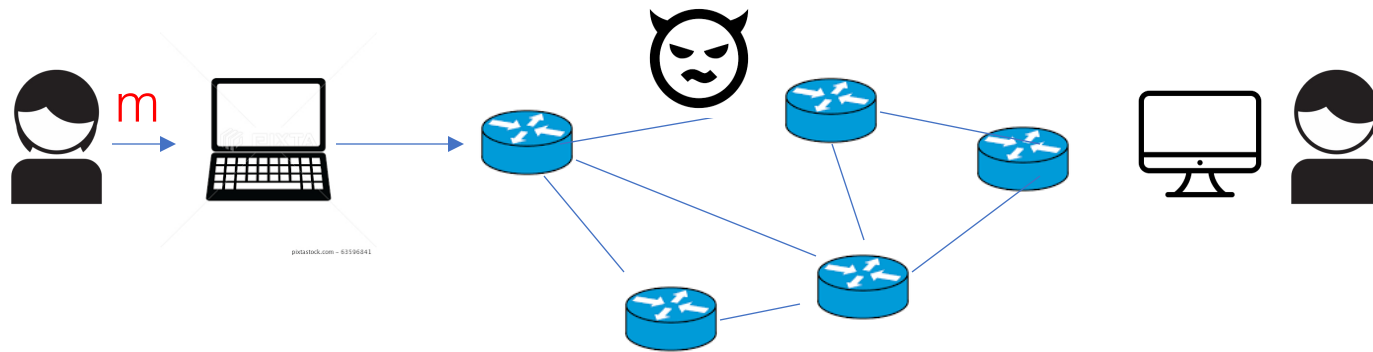
# Confidential communication over Internet



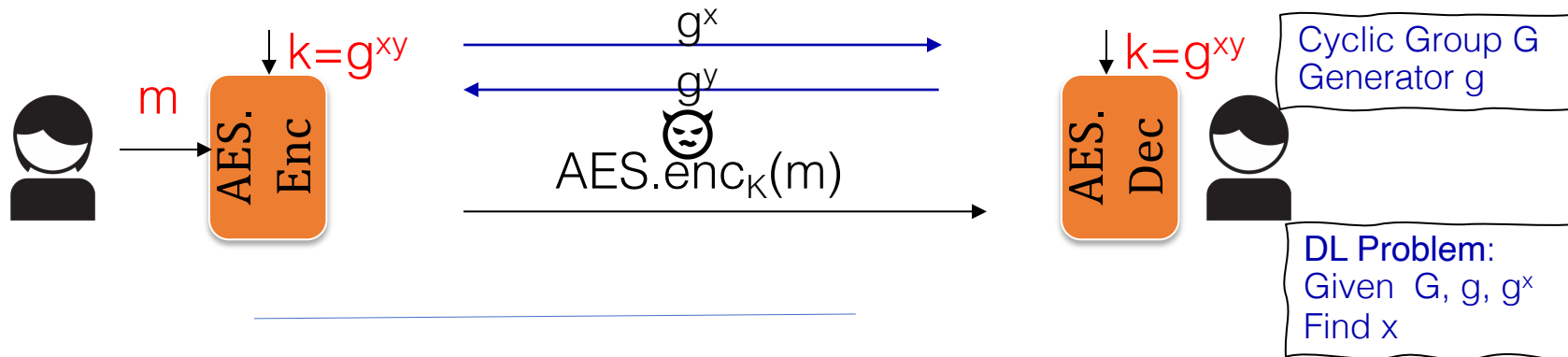
A secure key agreement (SKA) protocol establishes a key.



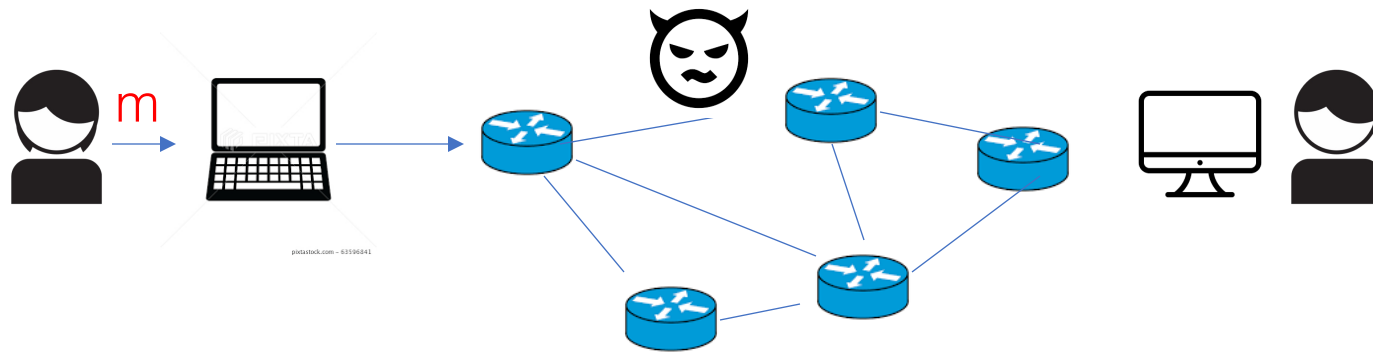
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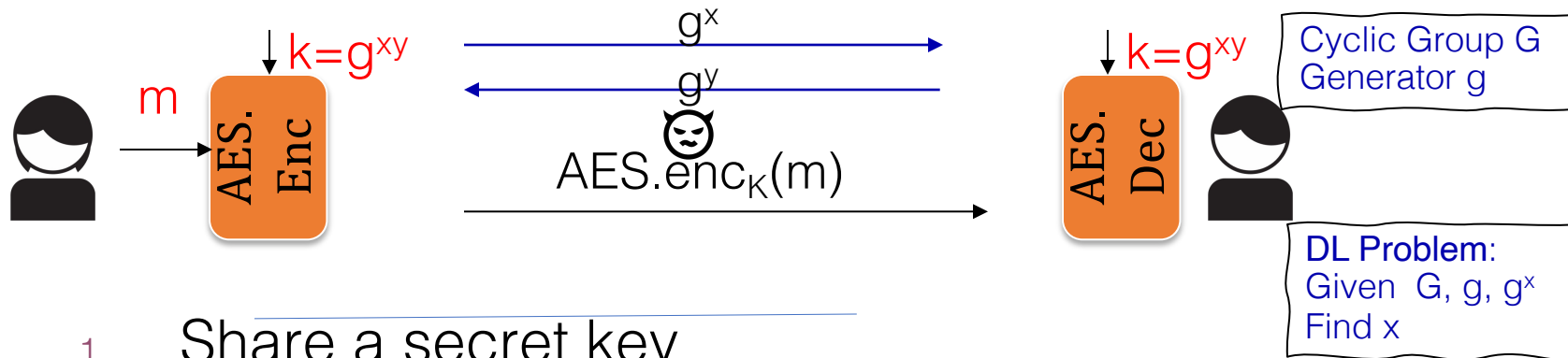
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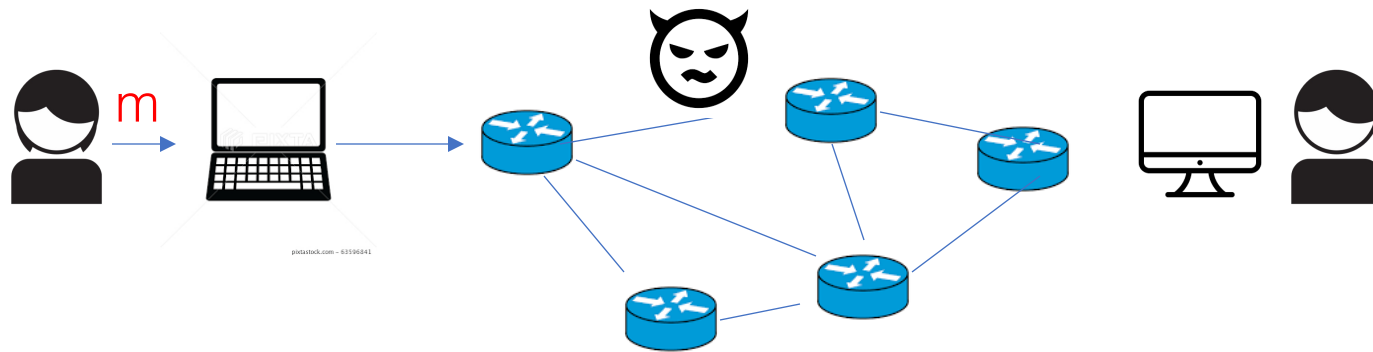


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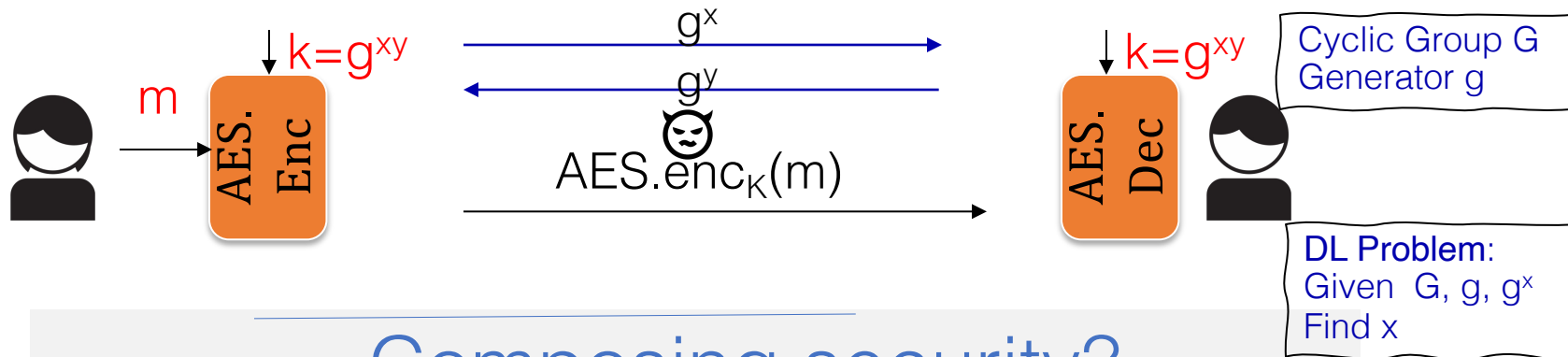


1. Share a secret key
2. Encrypt using a symmetric key cipher

# Confidential communication over Internet

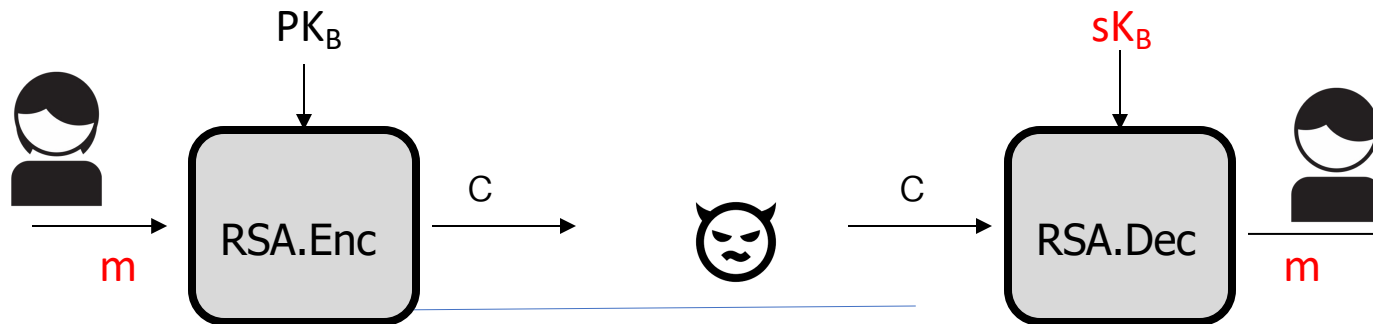
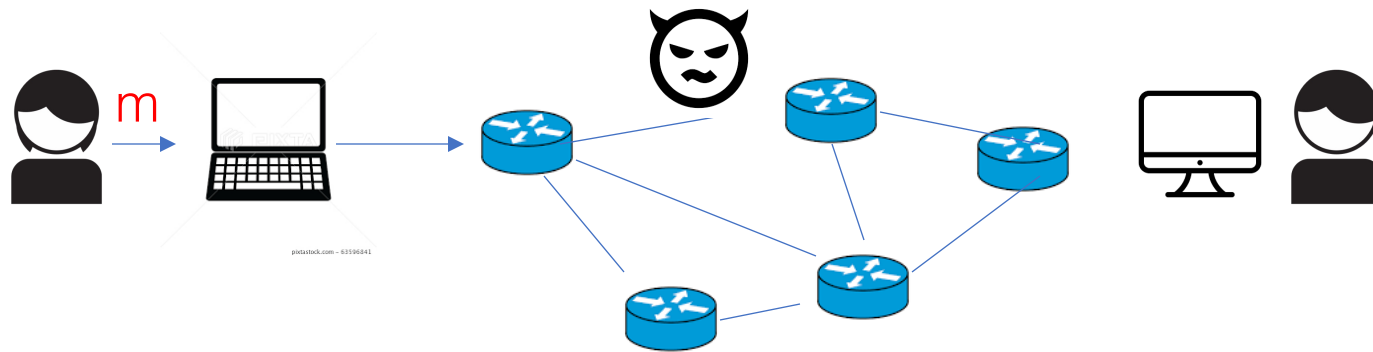


A secure key agreement (SKA) protocol establishes a key.



## Composing security?

# Confidential communication over Internet

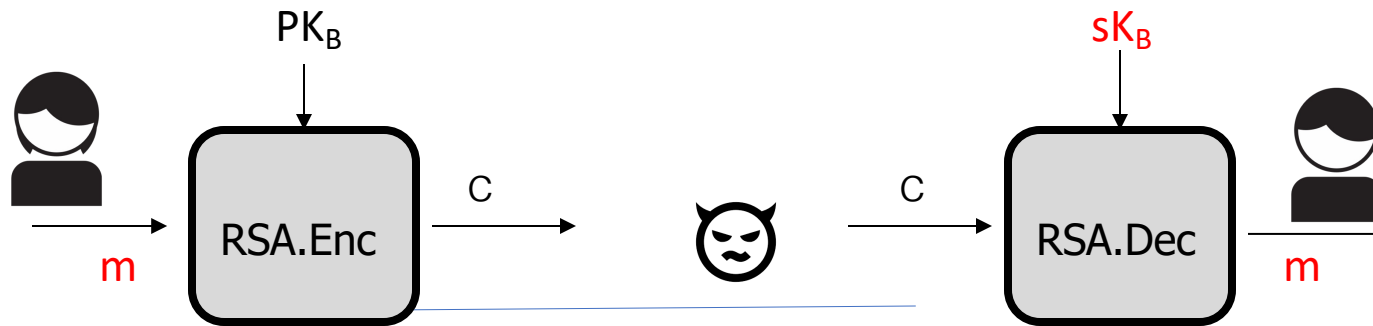
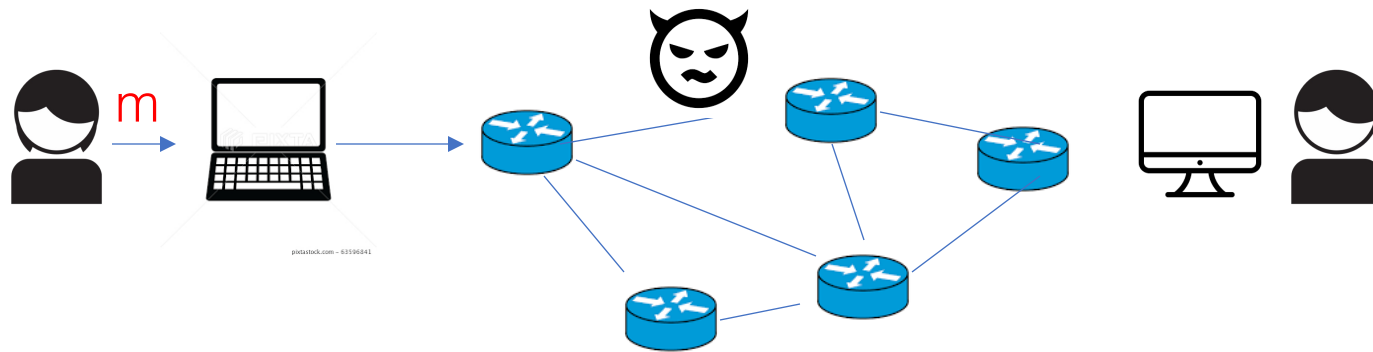


Public key encryption : No key sharing

ETSI 2023  
Computationally Inefficient



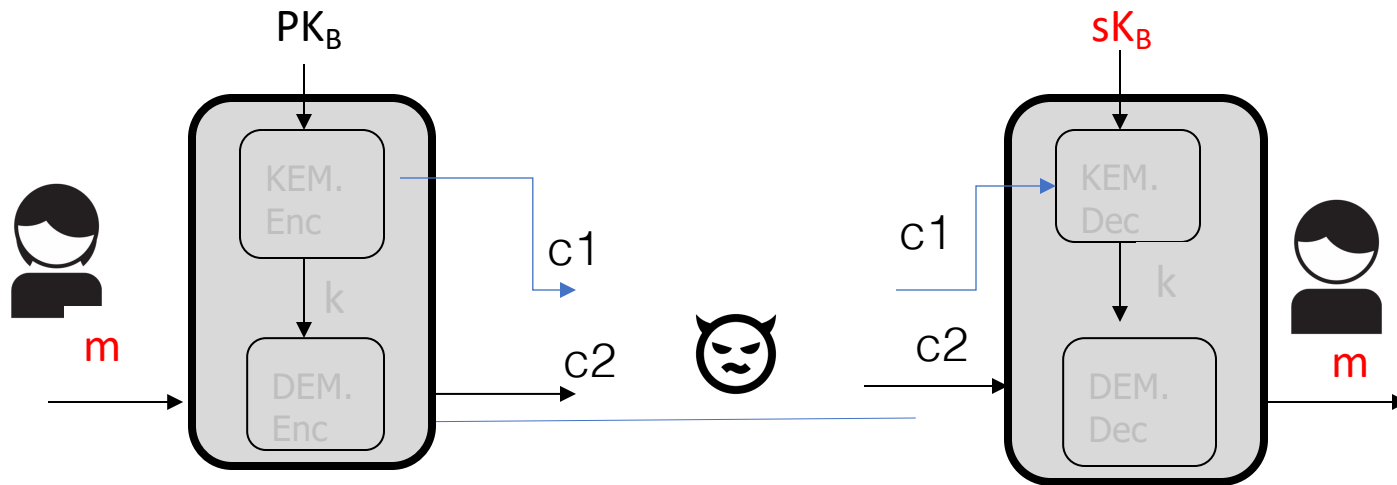
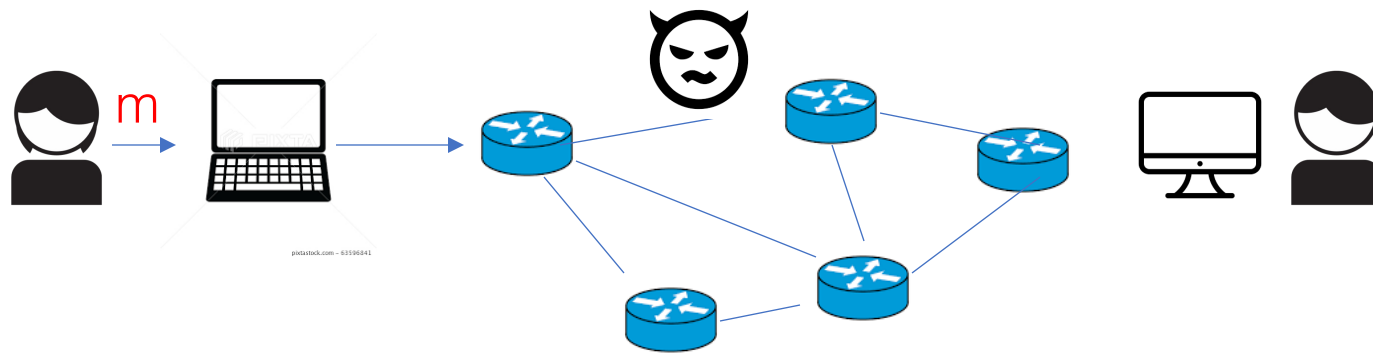
# Confidential communication over Internet



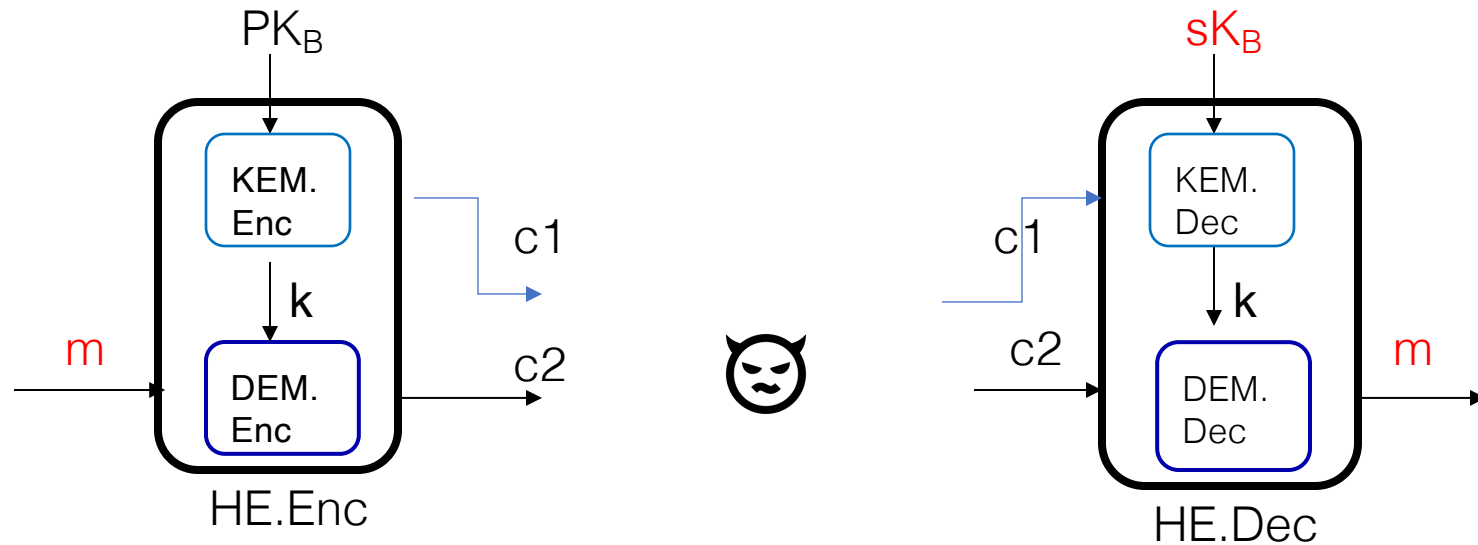
Public key encryption : No key sharing

Computationally Inefficient

# Hybrid Encryption: SKA + Encryption



# Hybrid Encryption: KEM + DEM



**KEM = (KEM.KGen, KEM.Enc, KEM.Dec)**

Key Generation:  $(PK, sK) \leftarrow \text{KEM.KGen}(1^\lambda)$

Key Encapsulation:  $(k, c_1) \leftarrow \text{KEM.Enc}(1^\lambda, PK)$

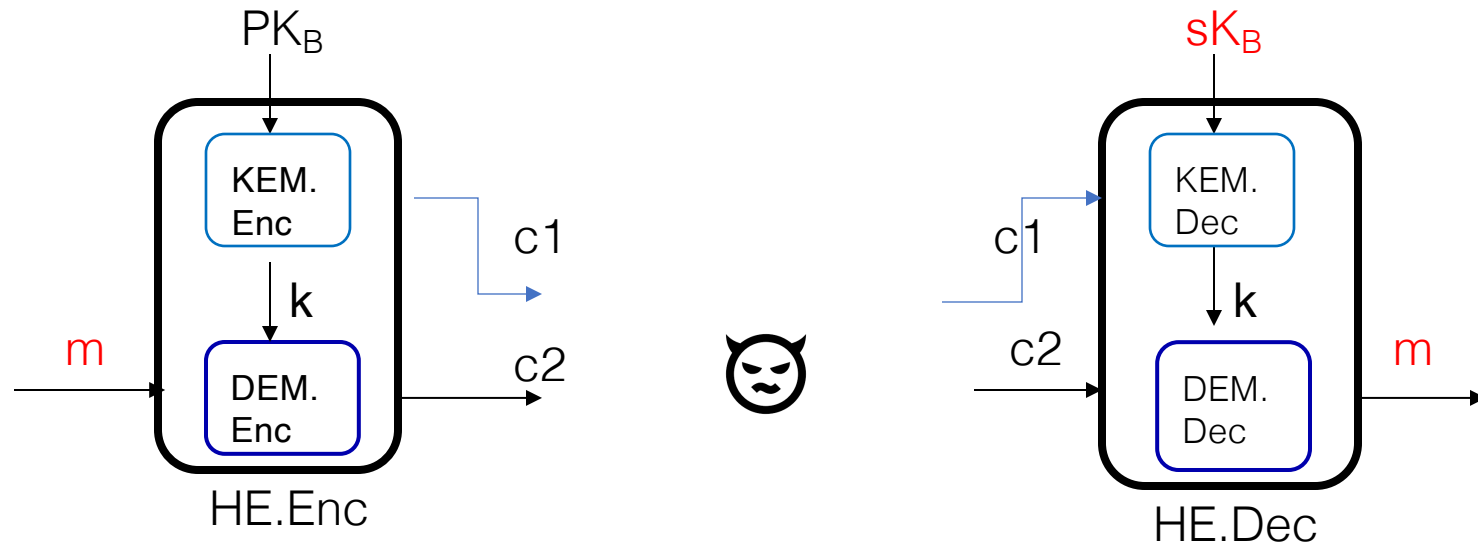
Key Decapsulation:  $k \leftarrow \text{KEM.Dec}(1^\lambda, sK, c_1)$

**DEM = (DEM.Enc, DEM.Dec)**

Data Encapsulation:  $c_2 \leftarrow \text{DEM.Enc}(1^\lambda, k, m)$

Data Decapsulation:  $m \leftarrow \text{DEM.Dec}(1^\lambda, k, c_1)$

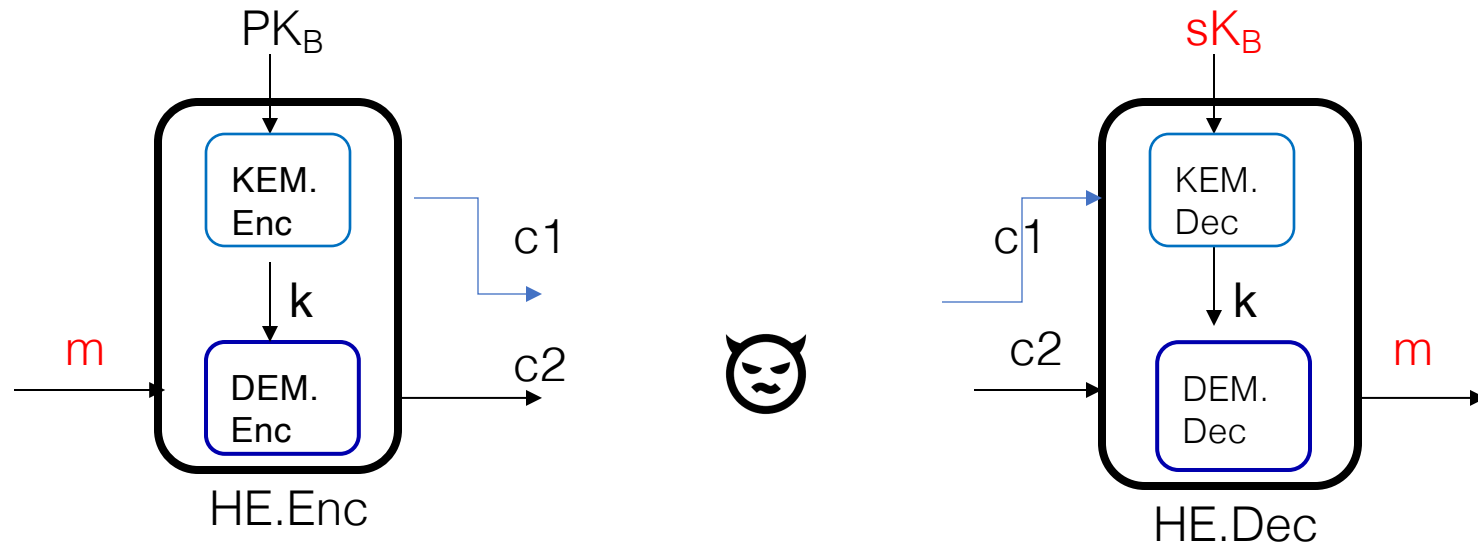
# Hybrid Encryption: KEM + DEM



## Security Theorem (Cramer & Shoup, Siam J of Comput, 2003)

If KEM and DEM are secure against adaptive chosen ciphertext attacks, then so is HPKE.

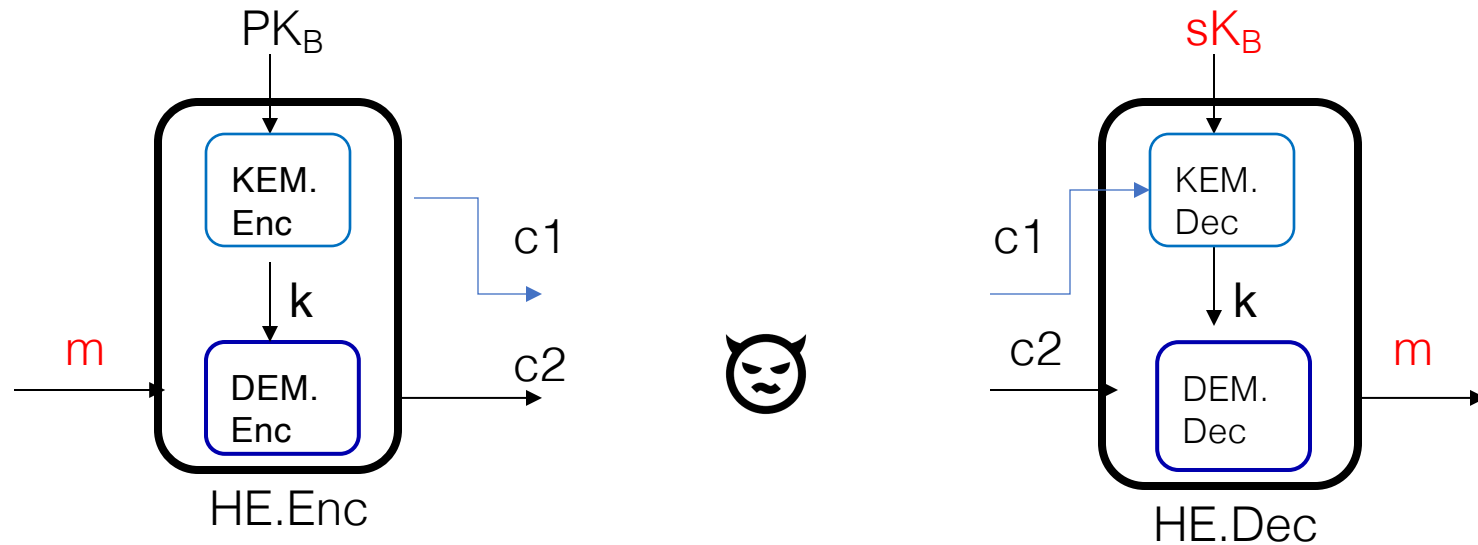
# Hybrid Encryption: KEM + DEM



## Kyber is a quantum-resistant IND-CCA2-secure KEM

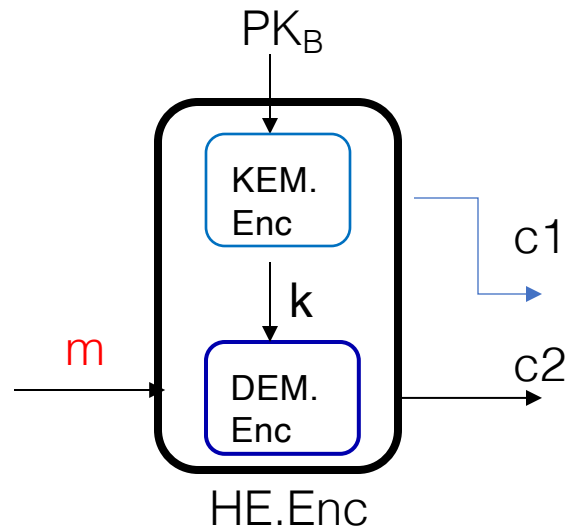
- Security is based on LWE problem over module lattices
- Standardized by NIST.
- Kyber is integrated into libraries and systems by industry.
  - Cloudflare integrated Kyber into CIRC;
  - Amazon supports hybrid modes involving Kyber AWS Key Management Service

# Hybrid Encryption: KEM + DEM



- Existing KEMs rely on computational assumptions
- Can we have KEM without computational assumptions?

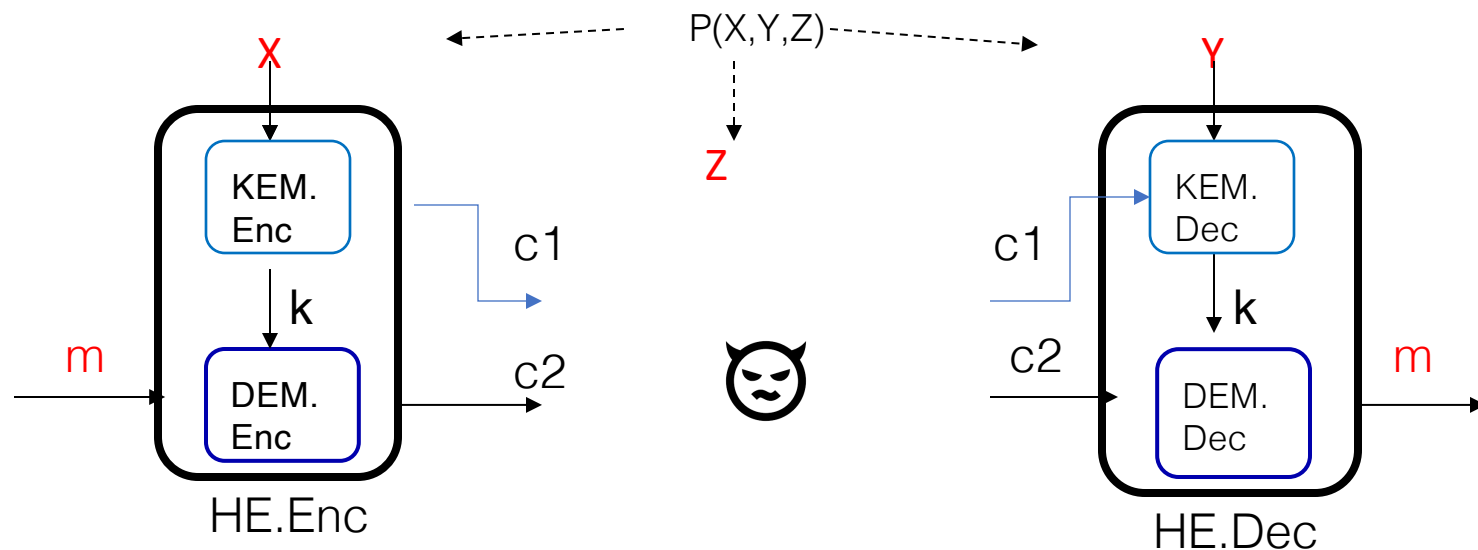
# HE without Computational Assumption



## Motivations:

- Computational assumptions remain vulnerable to advances in computing
- Expanding the crypto toolbox

# HE in Correlated Randomness Model



- KEM relies on correlated randomness
  - Information theoretic → iKEM
  - Computationally bounded → cKEM

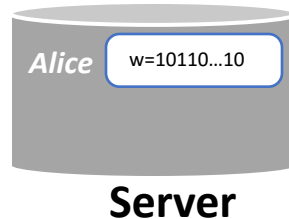


# Correlated Randomness Model:

## Where correlated randomness comes from?

- Biometrics

### 1. Registrations



### 2. Authentication

new scan



$w'=01110...10$

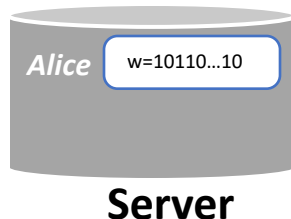
New scan code is correlated with the stored code.

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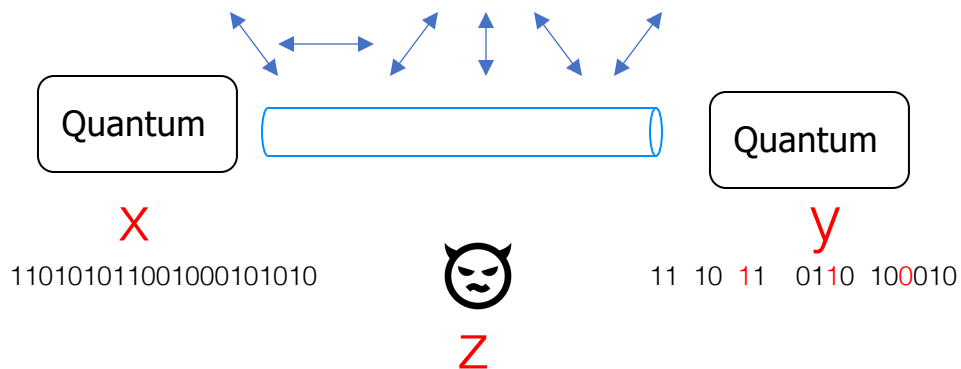
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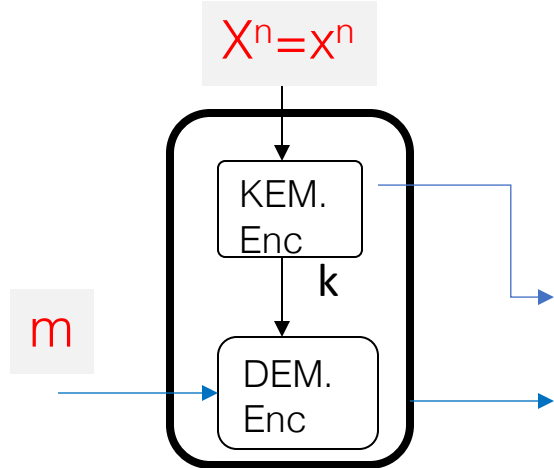
- Transmission over quantum channel



# Adversaries - Security

## Information theoretic

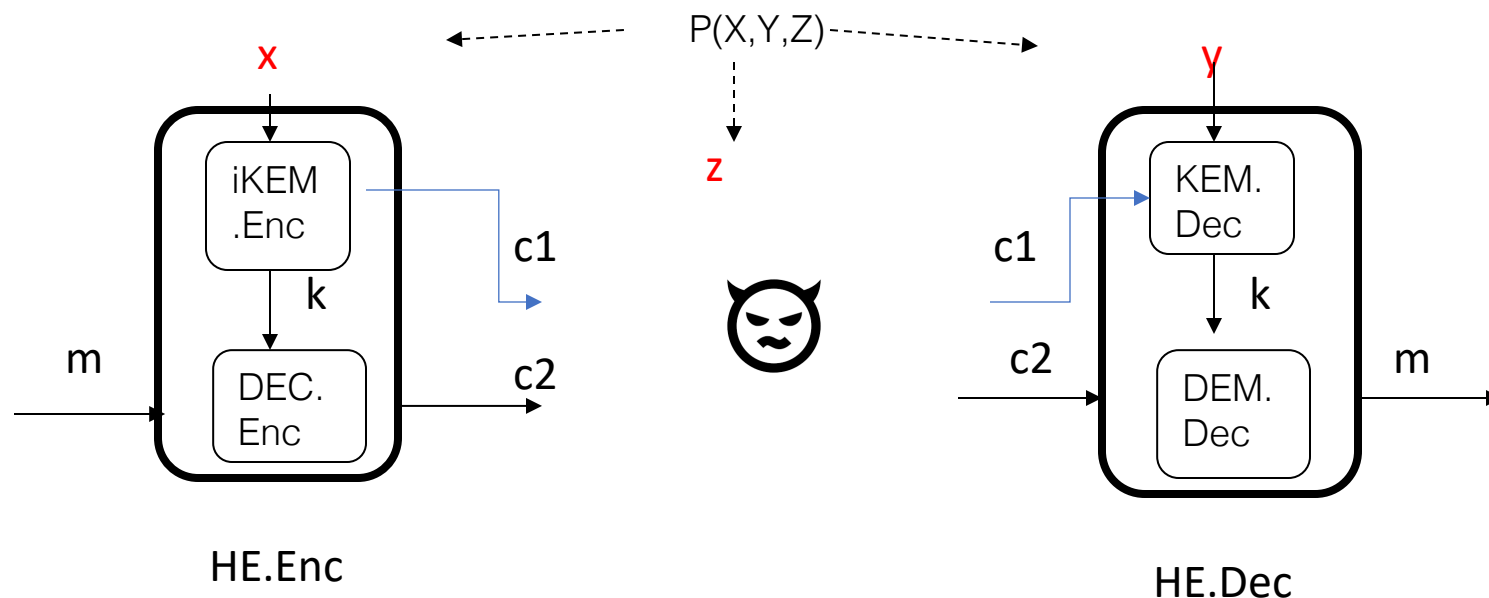
- Unbounded computation
- KEM & DEM are information theoretic
  - Long keys



## Computational Security

- Bounded computation
  - Access to quantum computer
- (iKEM or cKEM)+ DEM
  - Q-resistant security
    - Infeasible computation
      - Search a large key space
    - cKEM will use hard problems

# HE in Correlated Randomness model



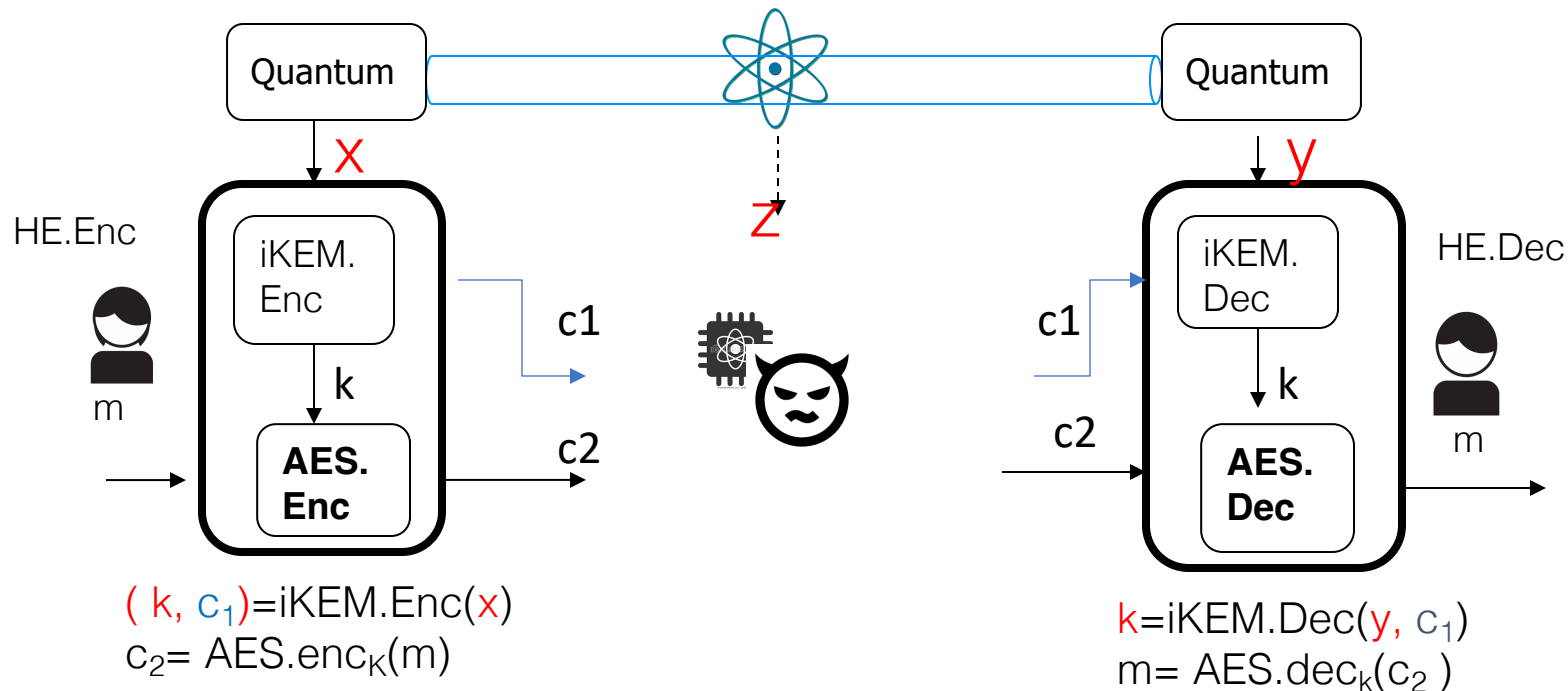
$(k, c_1) \leftarrow iKEM.Enc(1^\lambda, x)$   
 $c_2 \leftarrow DEM.Enc(1^\lambda, k, m)$   
 $(c_1, c_2) = HE.Enc(1^\lambda, x, m)$

$K \leftarrow iKEM.Dec(1^\lambda, c_1, y)$   
 $m \leftarrow DEM.Dec(1^\lambda, k, c_2)$

## Composition Theorem:

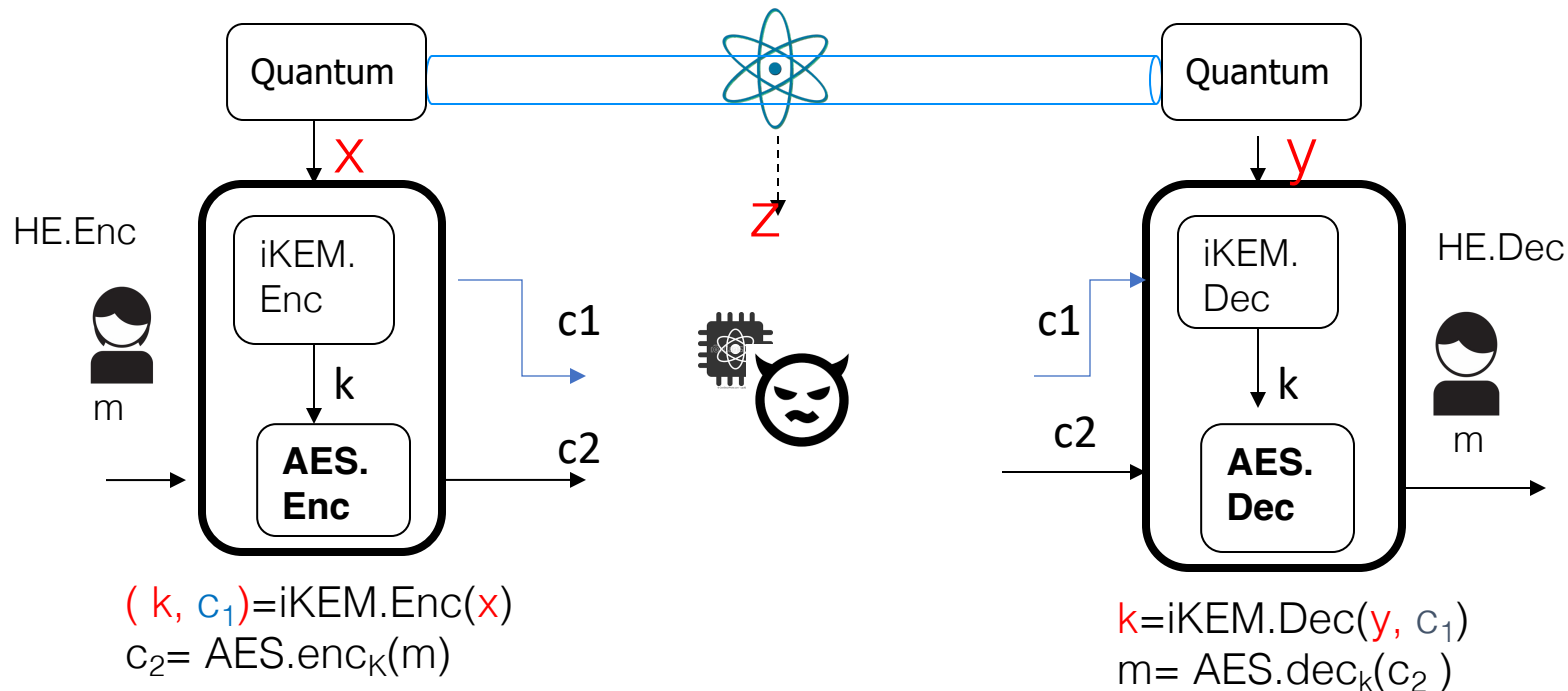
$IND-q-CEA\ iKEM + IND-OT\ DEM \Rightarrow IND-q-CPA\ HE$

# Quantum-enabled Quantum-resistant HE



- KEM relies on correlated randomness
  - Information theoretic  $\rightarrow$  iKEM
- Computational DEM
  - AES

# Quantum-enabled Quantum-resistant HE

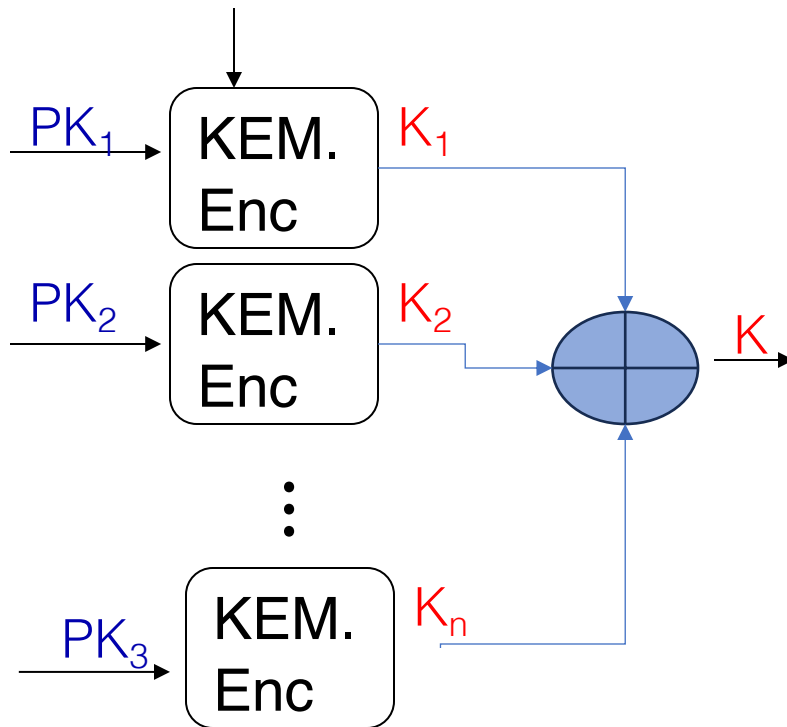


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Efficient Q-resistant hybrid encryption

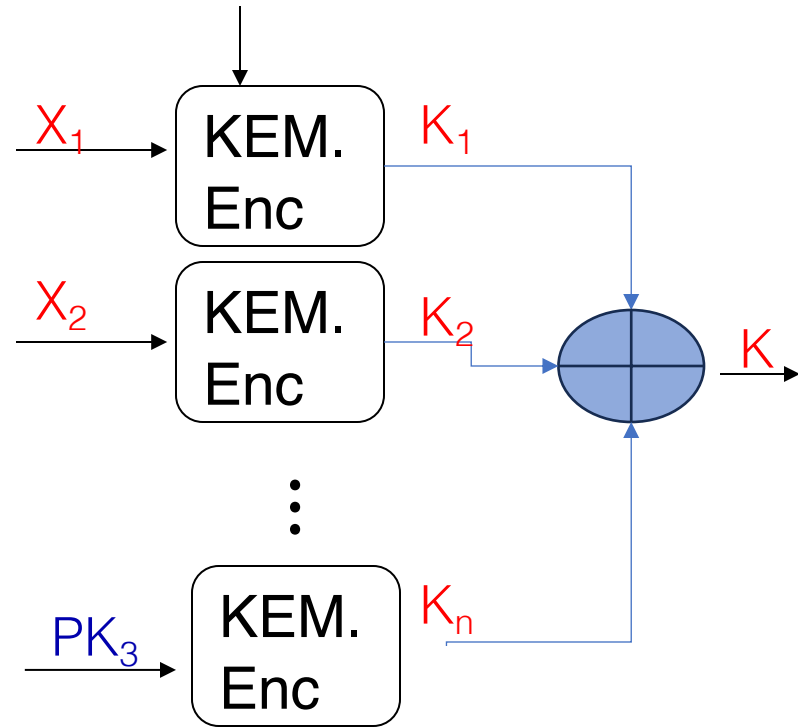
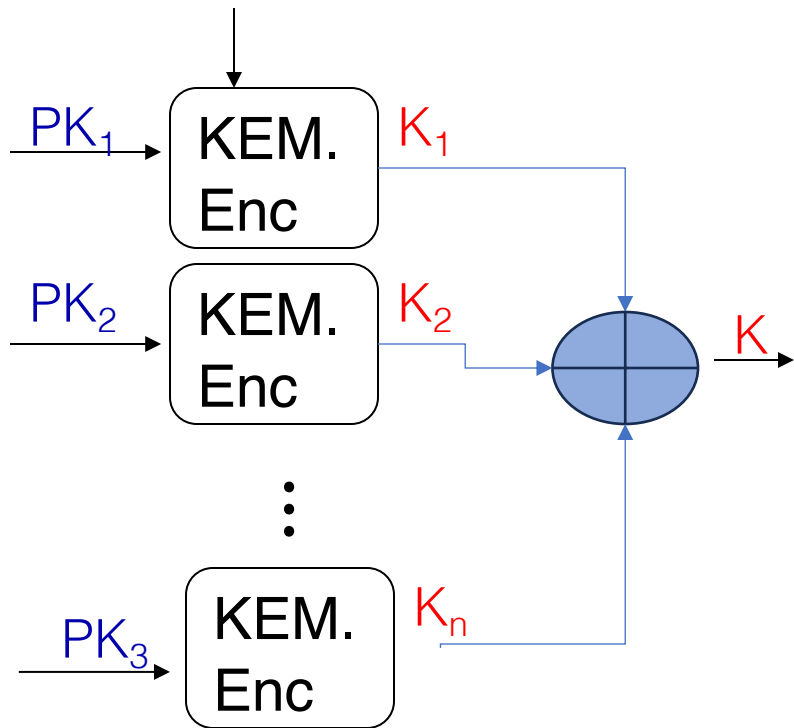
# KEM Combiners

- Cryptographic combiners



# KEM Combiners

- Cryptographic combiners





# Concluding remarks

- Quantum-resistant confidentiality
  - Proved security
  - **No** public key (hard problem)
  - Adversary with Access to quantum computer
  - Infeasible computation
    - Search a large key space
- Future work: quantum adversary
- Acknowledgement: Somnath Panja, Shaoquan Jiang, Setareh Sharifian, Nikita Tripathi, Ali Poostindouz

