

# **3GPP security hot topics:** LTE/SAE and Home (e)NB

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## Outline

- □ Some history and background
- □ SAE/LTE security: some highlights
- □ Home (e)NodeB security
- □ Summary



# Some history and background



## Some history (1/2)

For 3GPP Release 99 (frozen 2000), WG SA3 created
19 new specifications, e.g.

- > TS 33.102 "3G security; Security architecture"
- 5 specifications (out of these 19) originated by ETSI SAGE, e.g. TS 35.202 "KASUMI specification"
- For Release 4 (frozen 2001), SA3 was kept busy with GERAN security while ETSI SAGE originated again 5 new specifications, e.g.

**TS 35.205-208 for MILENAGE algorithm set** 

Release 5 (frozen 2002): SA3 added 3 new specifications, e.g.:

> TS 33.203 "IMS security"

> TS 33.210 "Network domain security: IP layer"



# Some history (2/2)

- Release 6 (frozen 2005): SA3 added 17 new specifications, e.g.:
  - > TS 33.246 "Security of MBMS"
  - > TS 33.220-222 "Generic Authentication Architecture"
- Release 7 (frozen 2007): SA3 added 13 new specifications
  - ETSI SAGE created 5 specifications for UEA2 & UIA2 (incl. SNOW 3G spec) (TS 35.215-218, TR 35.919)
- Release 8 (frozen 2008): SA3 has added 5 new specifications, e.g.:
  - > TS 33.401 "SAE: Security architecture"
  - > TS 33.402 "SAE: Security with non-3GPP accesses"
  - > (1-2 more TR's maybe still be included in Rel-8)



# SAE/LTE security: some highlights



## SAE/LTE: What and why?

SAE = System Architecture Evolution LTE = Long Term Evolution (of radio networks)

LTE offers higher data rates, up to 100 Mb/sec
SAE offers optimized (flat) IP-based architecture

**Technical terms:** 

- E-UTRAN = Evolved UTRAN (LTE radio network)
- > EPC = Evolved Packet Core (SAE core network)
- EPS = Evolved Packet System ( = RAN + EPC )



## **Implications on security**

□ Flat architecture:

- > All radio access protocols terminate in one node: eNB
- IP protocols also visible in eNB
- □ Security implications due to
  - > Architectural design decisions
  - Interworking with legacy and non-3GPP networks
  - Allowing eNB placement in untrusted locations
  - New business environments with less trusted networks involved
  - > Trying to keep security breaches as local as possible
- □ As a result (when compared to UTRAN/GERAN):
  - Extended Authentication and Key Agreement
  - More complex key hierarchy
  - More complex interworking security
  - Additional security for eNB (compared to NB/BTS/RNC)



## Authentication and key agreement



□ HSS generates authentication data and provides it to MME

Challenge-response authentication and key agreement procedure between MME and UE



## **Confidentiality and integrity of signalling**



- □ RRC signalling between UE and E-UTRAN
- □ NAS signalling between UE and MME
- □ S1 interface signalling
  - protection is not UE-specific
  - > optional to use



## **User plane confidentiality**



- □ S1-U protection is not UE-specific
  - (Enhanced) network domain security mechanisms (based on IPsec)
  - > Optional to use
- □ Integrity is not protected for various reasons, e.g.:
  - > performance
  - Iimited protection for application layer



## **Cryptographic network separation (1/2)**







## **Cryptographic network separation (2/2)**

- □ Authentication vectors are specific to the serving network
  - $\rightarrow$  AV's usable in UTRAN/GERAN cannot be used in EPS
- AV's usable for UTRAN/GERAN access cannot be used for E-UTRAN access
  - Solution by a "separation bit" in AMF field
- □ On the other hand, Rel-99 USIM is sufficient for EPS access
  - $\rightarrow$  ME has to check the "separation bit" (when accessing E-UTRAN)
- □ As one consequence, "EAP-AKA' " was created in IETF



## Handovers without MME involvement (1/2)





## Handovers without MME involvement (2/2)

- Handovers are possible directly between eNB's for performance reasons
- If keys would be passed as such, all eNB's in a "HO chain" would know all the keys → one compromised eNB would compromise all eNB's in the "HO chain"
- **Countermeasures:** 
  - > One-way function used before key is passed (*Backward security*)
  - MME is involved after the HO for further key passes (Forward security, effective after two hops)
  - When MME involved already during the HO, Forward security is effective already after one hop



# $\mathbf{K}_{\text{eNB}}$ derivations





## Interworking with UTRAN/GERAN (1/2)

- □ UE may be registered in both SGSN and MME simultaneously
  - → when moving from one system (source) to the other (target) both
  - cached keys (created earlier in the target system)
  - and
  - mapped keys (converted from the keys in the source system) may exist
  - > Note: cached keys only for Rel-8 SGSN, not for legacy SGSN



# Interworking with UTRAN/GERAN (2/2)

#### Idle mode transition

- From E-UTRAN to UTRAN: either mapped or cached keys are used (depending on the identity used in Routing Area Update Request)
- From UTRAN to E-UTRAN: cached keys are used but an exceptional case exists also

#### Handover

- From E-UTRAN to UTRAN: mapped keys are used
- From UTRAN to E-UTRAN: mapped keys are used but it is possible to activate the cached keys after HO completed (using key-change-onthe-fly procedure)



## Inter-working with non-3GPP networks (1/2)



Extract from TS 23.402 (one of several architecture figures)



## Inter-working with non-3GPP networks (2/2)

- Three options for mobility between 3GPP and non-3GPP networks:
  - Proxy Mobile IP: no user-specific security associations between the Proxy and Home Agent
  - > Client MIPv4: tailor-made security mechanisms are used
  - > Dual Stack MIPv6: IPsec with IKEv2 is used between UE and HA
- IPsec tunnel (with evolved Packet Data Gateway) is used in case the non-3GPP network is untrusted by the operator (of EPS network)
- Authentication is run by EAP-AKA or EAP-AKA' procedures, in both cases based on USIM





# Home (e) Node B security



# Home (e)NB architecture



Figure from draft TR 33.820

One of the key concepts: Closed Subscriber Group

Note: Rest of the talk: Home (e)NB denoted by HeNB



## **Threats**

**Compromise of HeNB credentials** e.g. cloning of credentials Physical attacks on HeNB > e.g. physical tampering Configuration attacks on HeNB > e.g. fraudulent software updates Protocol attacks on HeNB > e.g. man-in-the-middle attacks □ Attacks against the core network > e.g. Denial of service Attacks against user data and identity privacy > e.g. by eavesdropping

Attacks against radio resources and management



## **Several sources of Security Requirements**

- □ (Additional) requirements for eNB due to SAE/LTE security architecture (TS 33.401)
- Requirements stemming from threats due to home placement (TR 33.820)
- □ Requirements due to *Closed Subscriber Group* concept



## Countermeasures

- Mutual authentication between the HeNB and the (rest of) network
- Security tunnel establishment for backhaul link
- Trusted Environment inside HeNB
  - ➢ e.g. secure execution
- □ Access Control mechanisms (for Closed Subscriber Groups)
- □ Security mechanisms for OAM
- **Hosting party authentication (if used) with Hosting Party Module**
- **d** etc..





# Summary

#### □ SAE/LTE security

- New architecture and business environment require enhancements to 3G security
- > Radio interface user plane security terminates in base station site
- Cryptographic separation of keys
- Forward/backward security in handovers
- Different security mechanisms in many inter-working cases with both 3GPP and non-3GPP access networks

#### □ Home (e)NB security

- > New architecture with more exposed locations of NB's
- > New types of threats
- Many new countermeasures needed



# For more information: www.3gpp.org